

Syntax Analysis

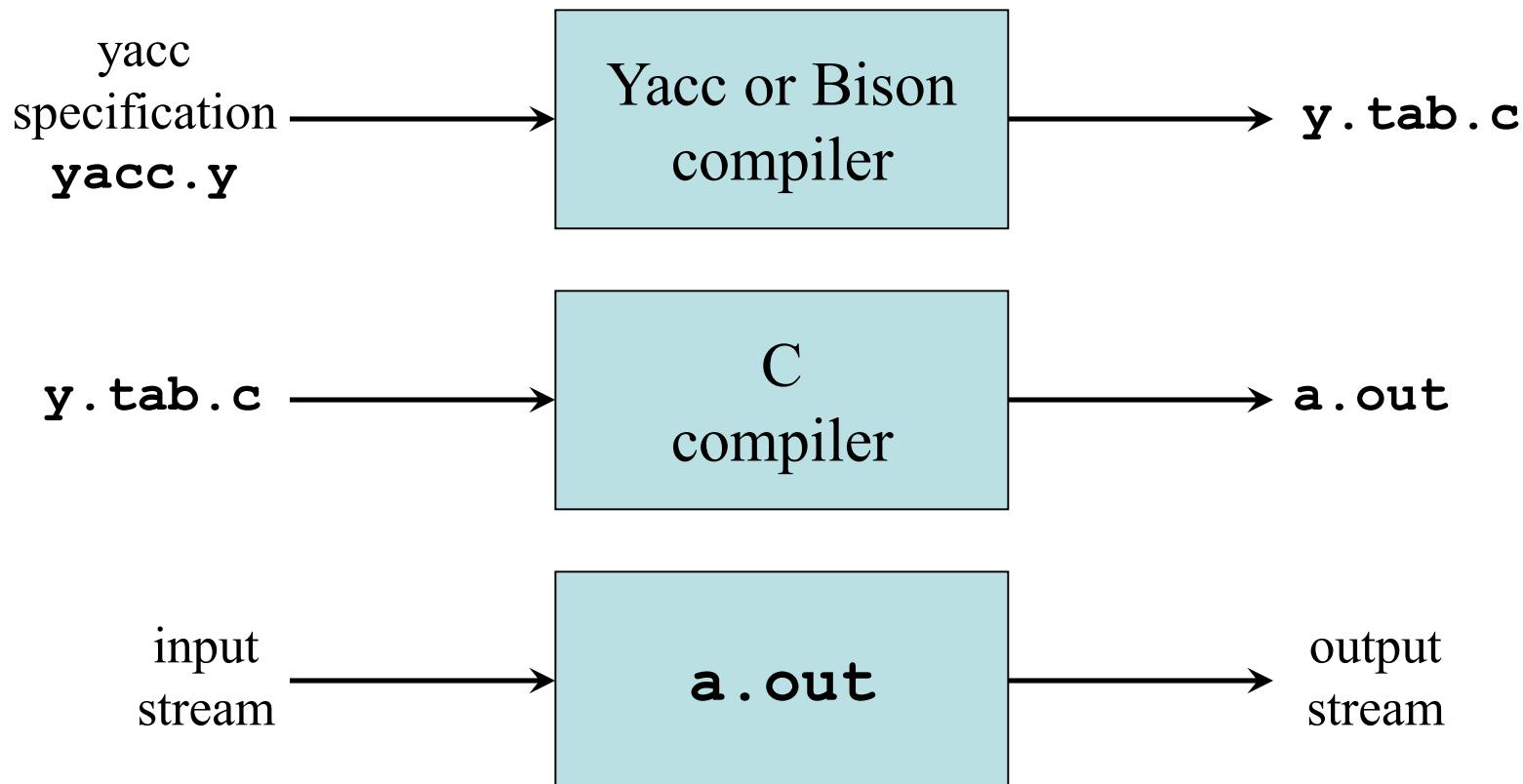
Part III

Chapter 4

ANTLR, Yacc, and Bison

- *ANTLR* tool generates LL(k) parsers
- *Yacc* (Yet Another Compiler Compiler) generates LALR(1) parsers
- *Bison* (Yacc improved)

Creating an LALR(1) Parser with Yacc/Bison



Yacc Specification

- A *yacc specification* consists of three parts:
yacc declarations, and C declarations in % { % }
%%
translation rules
%%
user-defined auxiliary procedures
- *Translation rules* are grammar productions and actions:
*production*₁ { *semantic action*₁ }
*production*₂ { *semantic action*₂ }
...
*production*_{*n*} { *semantic action*_{*n*} }

Writing a Grammar in Yacc

- Productions in Yacc are of the form
 - $Nonterminal : \text{tokens/nonterminals} \{ action \}$
 - $| \text{tokens/nonterminals} \{ action \}$
 - ...
 - ;
- Tokens that are single characters can be used directly within productions, e.g. `+'
 - Named tokens must be declared first in the declaration part using
%token TokenName

Synthesized Attributes

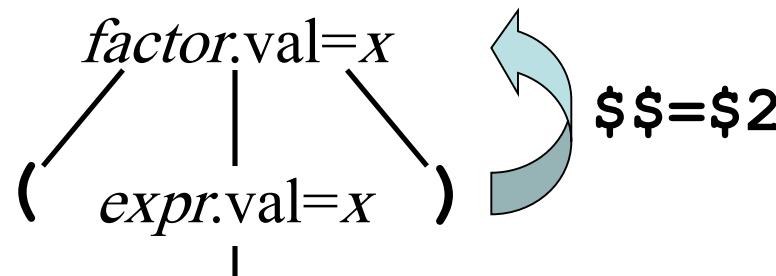
- Semantic actions may refer to values of the *synthesized attributes* of terminals and nonterminals in a production:

$$X: Y_1 \ Y_2 \ Y_3 \dots \ Y_n \quad \{ \text{action} \}$$

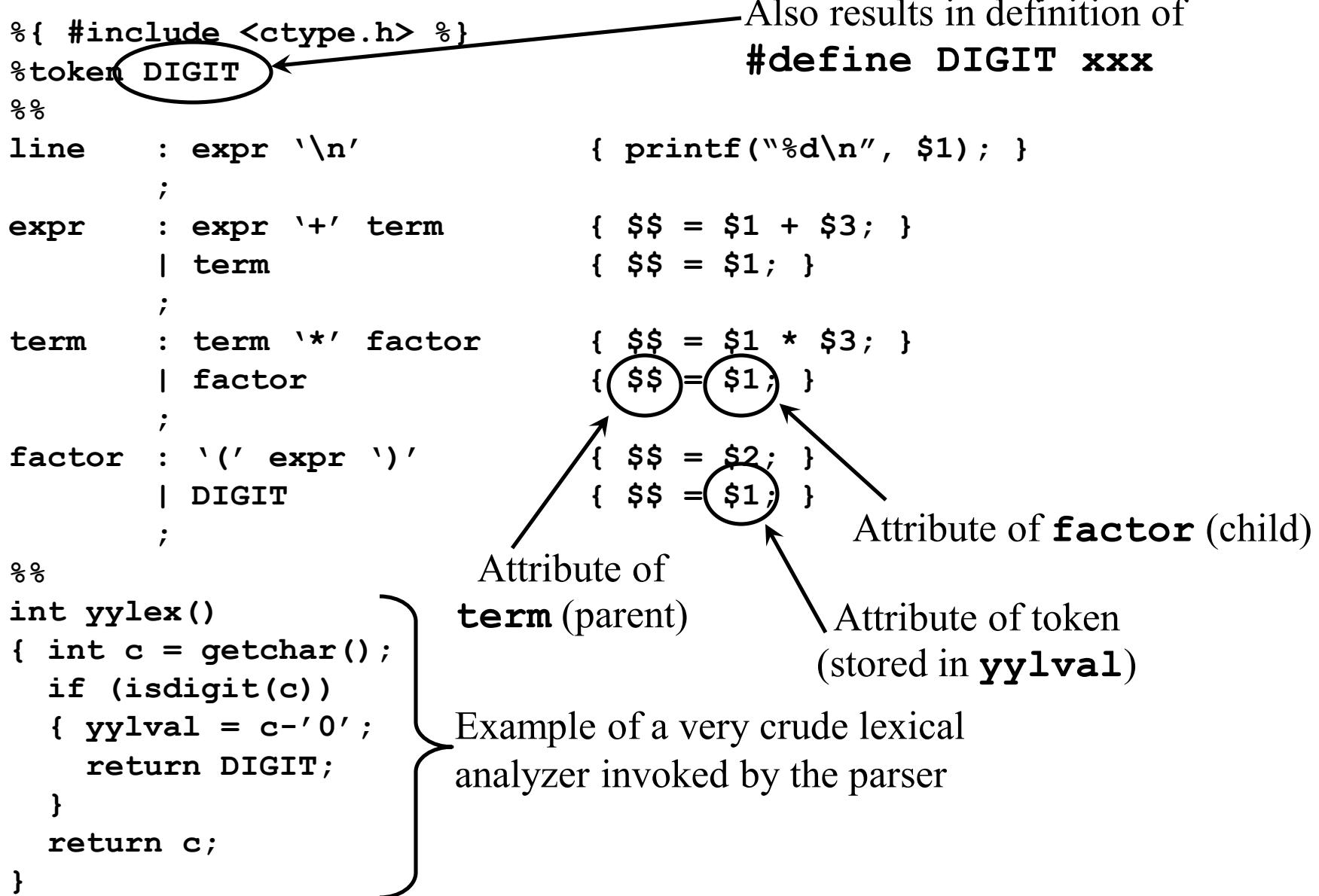
- $\$\$$ refers to the value of the attribute of X
- $\$i$ refers to the value of the attribute of Y_i

- For example

```
factor : ' (' expr ') ' { $$=$2 ; }
```



Example 1



Dealing With Ambiguous Grammars

- By defining operator precedence levels and left/right associativity of the operators, we can specify ambiguous grammars in Yacc, such as
$$E \rightarrow E+E \mid E-E \mid E^*E \mid E/E \mid (E) \mid -E \mid \mathbf{num}$$
- To define precedence levels and associativity in Yacc's declaration part:

```
%left '+' '-'
%left '*' '/'
%right UMINUS
```

Example 2

```
%{
#include <ctype.h>
#include <stdio.h>
#define YYSTYPE double
%}

%token NUMBER
%left '+'
%left '-'
%left '*'
%left '/'
%right UMINUS
%%

lines : lines expr '\n'          { printf("%g\n", $2); }
      | lines '\n'
      | /* empty */
;

expr  : expr '+' expr          { $$ = $1 + $3; }
      | expr '-' expr          { $$ = $1 - $3; }
      | expr '*' expr          { $$ = $1 * $3; }
      | expr '/' expr          { $$ = $1 / $3; }
      | '(' expr ')'
      | '-' expr %prec UMINUS { $$ = -$2; }
      | NUMBER
;

%%
```

Double type for attributes
and **yyval**

Example 2 (cont'd)

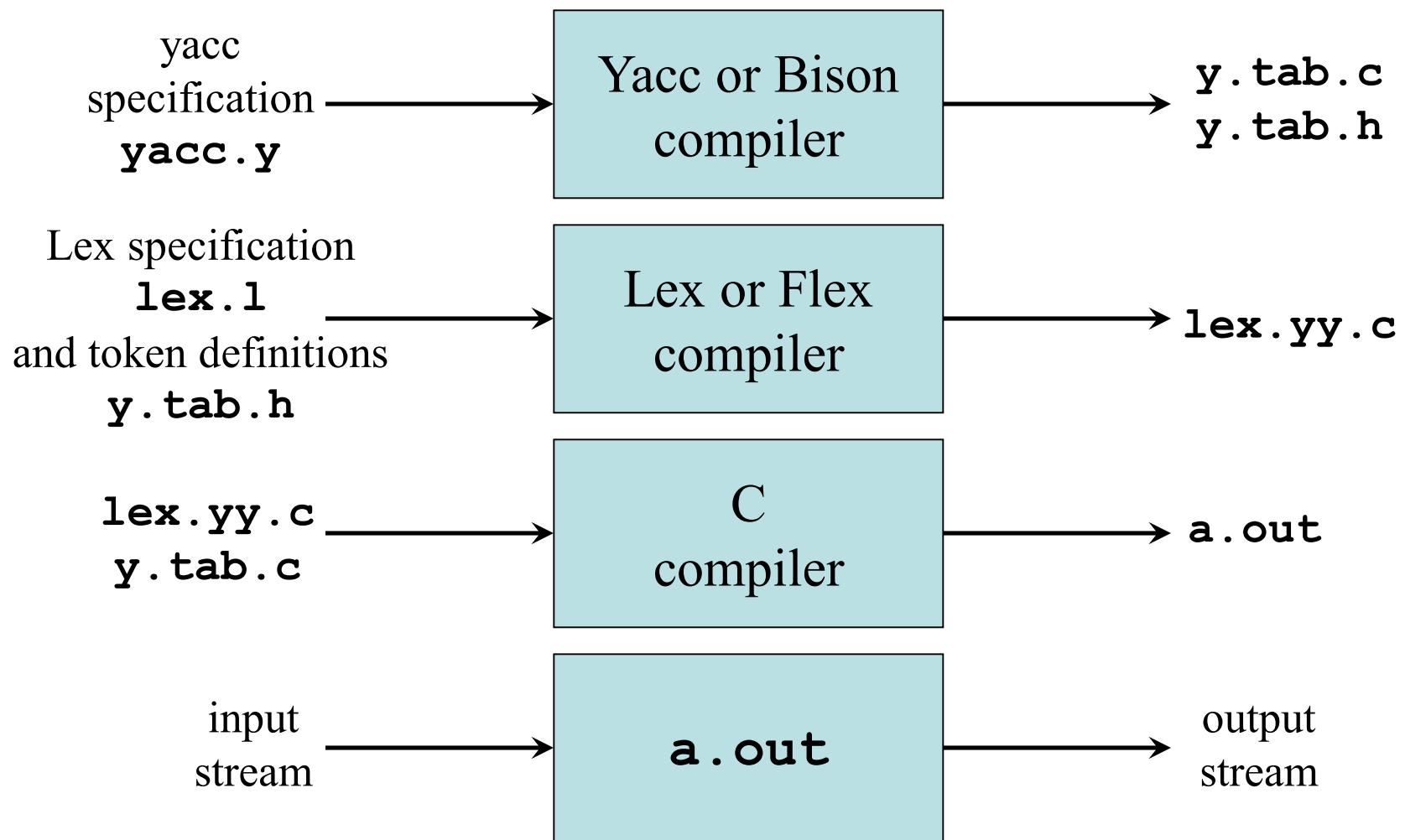
```
%%
int yylex()
{ int c;
  while ((c = getchar()) == ' ')
    ;
  if ((c == '.') || isdigit(c))
  { ungetc(c, stdin);
    scanf("%lf", &yyval);
    return NUMBER;
  }
  return c;
}
int main()
{ if (yparse() != 0)
  fprintf(stderr, "Abnormal exit\n");
  return 0;
}
int yyerror(char *s)
{ fprintf(stderr, "Error: %s\n", s);
}
```

Crude lexical analyzer for
fp doubles and arithmetic
operators

Run the parser

Invoked by parser
to report parse errors

Combining Lex/Flex with Yacc/Bison



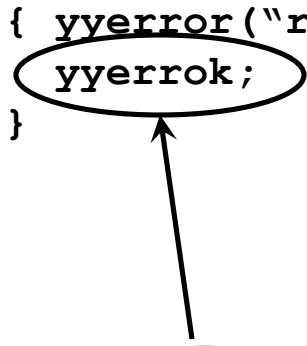
Lex Specification for Example 2

```
%option noyywrap
%{
#include "y.tab.h"           ← Generated by Yacc, contains
extern double yyval;         ← #define NUMBER xxx
%
number [0-9]+\.?|[0-9]*\. [0-9]+
%%
[ ]                         { /* skip blanks */ }
{number}        { sscanf(yytext, "%lf", &yyval);
                  return NUMBER;
}
\n|.           { return yytext[0]; }
```

```
yacc -d example2.y
lex example2.l
gcc y.tab.c lex.yy.c
./a.out
```

```
bison -d -y example2.y
flex example2.l
gcc y.tab.c lex.yy.c
./a.out
```

Error Recovery in Yacc

```
%{  
...  
%}  
...  
%%  
lines : lines expr '\n' { printf("%g\n", $2; ) }  
| lines '\n'  
| /* empty */  
| error '\n'  
;  
...  
  
Error production:  
set error mode and  
skip input until newline  
{ yyerror("reenter last line: ");  
  yyerrok;  
}
```

Reset parser to normal mode